

## Assignment no. 4

due: May 23rd, 2022

**Exercise 4.1** Prove that the number of inner nodes of the search structure  $\mathcal{D}$  of algorithm TRAPEZOIDALMAP increases by  $k_i - 1$  in iteration  $i$ , where  $k_i$  is the number of new trapezoids in the trapezoidal map  $\mathcal{T}(S_i)$ , and hence the number of new leaves of  $\mathcal{D}$ . (The exercise uses the notation of the book CGAA.)

**Exercise 4.2** A polygon  $P$  is called star-shaped if a point  $p$  in the interior of  $P$  exists such that, for any other point  $q$  in  $P$ , the line segment  $pq$  lies in  $P$ .

(a) Given a simple polygon  $P$  with  $n$  vertices, describe a procedure running in expected  $O(n)$  time, to determine whether  $P$  is star-shaped. Prove the correctness of the approach.

(b) Given a star-shaped polygon  $Q$  with  $n$  vertices, show that after expected  $O(n)$  preprocessing time, one can determine whether a query point lies in  $Q$  in worst-case  $O(\log n)$  time.

**Exercise 4.3** Design an algorithm with running time  $O(n \log n)$  for the following problem: Given a set  $P$  of  $n$  points, determine a value of  $\varepsilon > 0$  such that the shear transformation  $\Phi : (x, y) \rightarrow (x + \varepsilon y, y)$  (see Section 6.3 in CGAA) does not change the order, in  $x$ -direction, of points with unequal  $x$ -coordinates.

---

**Exercise 4.4** Let  $L$  be a set of  $n$  lines in the plane. Give an  $O(n \log n)$  time algorithm to compute an axis-parallel rectangle that contains all the vertices of the arrangement  $\mathcal{A}(L)$  in its interior.

**Exercise 4.5** Hopcroft's problem is to decide, given  $n$  lines and  $n$  points in the plane, whether any point is contained in any line. Give an  $O(n^{3/2} \log n)$  time algorithm to solve Hopcroft's problem. Hint: Give an  $O(n \log n)$  time algorithm to decide, given  $\sqrt{n}$  lines and  $n$  points in the plane, whether any point is contained in any line.

**Exercise 4.6** Let  $S$  be a set of  $n$  segments in the plane. A line  $\ell$  that intersects all the segments of  $S$  is called a transversal or stabber for  $S$ .

(a) Give an  $O(n^2)$  algorithm to decide if a stabber exists for  $S$ .

(b) Now assume that all the segments in  $S$  are vertical. Give a randomized algorithm with  $O(n)$  expected running time that decides if a stabber exists for  $S$ . (CGAA Ex. 8.16)